

RAT Test Run

Question 1: In TCP, the Sequence Number in a segment refers to:

- A The total number of bytes sent so far in the connection.
- B The number of bytes currently unacknowledged.
- C The byte number of the first data byte carried in this segment (relative to the connection's initial sequence number).
- D The byte number of the last data byte carried in this segment.

Question 2: Why does TCP use an exponential weighted moving average for RTT estimation? Why not a simple arithmetic average?

- A Because with an **arithmetic average** newer samples would have higher influence and reflect newer changes in the link.
- B Because with an **exponential weighted moving average** there would be no difference between older and newer samples.
- C Because with an **arithmetic average** older and newer samples would have the same influence.
- D Because with an **exponential weighted moving average** older samples would have higher influence and better reflect the overall history of the link.

Question 3: TCP uses the RTT estimate to compute what value?

- A To compute the fast retransmit threshold.
- B To compute the sender window.
- C To compute the retransmission timeout value.
- D To compute the receive window.

Question 4: Imagine you want to download a file with a size of 1000 MiB. Assume a Maximum Segment Size (MSS) of 1024 B, a total of 64 B of transport, network and data-link header added to each segment, and a constant bandwidth of 65536 Kbit/s. For simplicity, ignore flow and congestion control, i.e., the sender can pump out the segments back to back and continuously. **How long does it take to transmit the file?** (Hints: It's easier to work with powers of 2; $65536 = 2^{16}$; MiB \neq Mbit)

- A It takes 128.125 s (i.e., $2^{-3} + 2^7$).
- B It takes 128 s (i.e., 2^7).
- C It takes 136 s (i.e., $2^3 + 2^7$).
- D It takes 17 s (i.e., $2^0 + 2^4$).

Question 5: Suppose *Host A* sends two TCP segments back-to-back to *Host B* over a TCP connection. The first segment has the sequence number 5212; the second has the sequence number 5812. How much data (in bytes) is in the **first segment**?

- A 600 Bytes
- B 580 Bytes assuming a TCP header of 20 B (i.e., without options).
- C Can't say because it depends on the MSS value.
- D 599 Bytes

Question 6: Suppose *Host A* sends two TCP segments back-to-back to *Host B* over a TCP connection. The first segment has the sequence number 5212; the second has the sequence number 5812. Assuming the first segment is lost but the second segment arrives at *Host B*, what will be the **acknowledgment number** sent from *Host B* to *Host A*?

- A 5211
- B 5813
- C 5212
- D 5812

Question 7: Suppose *Host A* sends two TCP segments back-to-back to *Host B* over a TCP connection. The first segment has the sequence number 5212; the second has the sequence number 5812. Assuming the ACK for the first segment is lost but the ACK for the second segment arrives at *Host A* before any timeout, **what will *Host A* send next?**

- A A retransmission of the TCP segment with the sequence number 5812.
- B A new TCP segment with a new sequence number > 5812 .
- C A retransmission of the TCP segment with the sequence number 5212.
- D A new TCP segment with a the sequence number 5812.

Question 8: A TCP fast retransmit occurs:

- A After a TCP timeout timer expires for a given segment.
- B When three duplicate ACKs are received by the sender.
- C When one duplicate ACK is received by the sender.
- D If the TCP timeout timer is set with a very short value, resulting in faster retransmissions.

Question 9: TCP's flow control depends on:

- A the network's window that is included in the IP datagrams sent over the network.
- B the sender's window that is included in the TCP segments sent to the receiver.
- C the last sent byte being lower than the receive window, **regardless** of the last ACK'ed byte.
- D the receiver's window that is included in the TCP segments sent by the receiver.

Question 10: UDP's flow control depends on:

- A** the sender's window that can optionally be included in the UDP datagrams sent to the receiver.
- B** the network's window that is included in the IP datagrams sent over the network.
- C** nothing because UDP does not offer flow control, which can only be supported by the transport layer.
- D** UDP has no built-in flow control; any pacing is purely application-defined.